Dynamic Position Auctions with Consumer Search

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Dynamic Position Auctions

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Overview

Road Map

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Position Auctions

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• Position Auctions (without consumer search)

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Position Auctions (without consumer search) Motivation

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- Static Model

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- Welfare implications not well-understood

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 - By convention: $\theta_j = 0$

 $(i = M + 1, \ldots, N)$

Generalized Second-Price (GSP) Auction

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Payoff to
$$\pi_j: heta_j(m{q}_{\pi_j}-m{b}_{\pi_{j+1}})$$

Equilibria

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- Advertisers may update bids at any time
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- Do equilibrium results break down once dynamics are introduced?

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Cary et al. (2008) Dynamic Model

• Extends Edelman, Ostrovsky, and Schwarz (2007)

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 - Advertisers play a "best-response" strategy

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Unique fixed point: Bids follow a recursive formula,

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- If all advertisers play the Restricted Balanced Bidding strategy, then their bids converge to the fixed point; this convergence is efficient.
- The dynamic model is "well-approximated" by the static model.

Road Map

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- Motivation
- Static Model (Edelman, Ostrovsky, and Schwarz (2007))
- Dynamic Model (Cary et al. (2008))
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Our Dynamic Model

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Framework & Conventions

- Extends Athey and Ellison (2008)
- Ovnamic setting
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 - Consumers ignorant of dynamics

Framework & Conventions

Balanced Bidding

. . .

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 - Athey and Ellison (2008) Envy-Free Equilibrium

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Our Model

Framework & Conventions

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Main Result

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• The dynamic model is "well-approximated" by the static model.

Parameters

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Lemma

At every round $t > t_1 = 2 + \log_{\gamma^{**}}((1 - \gamma^{**})(q_M - q_{M+1})/q_{M+1})$:

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21 / 31

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Proof Approach. "Dynamical system in γ^{**} ."

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Convergence of the *M* Positions • By the Lemma

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 - If $\pi(P) = \{1, \dots, M\}$, then we are done.

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- (Significant divergence from Cary et al. (2008))

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Lemma

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Lemma

Let $\epsilon = \frac{G(\bar{q}_M)}{2G(\bar{q}_1)} (1 - \gamma^{**}) \min_{\phi \neq \phi'} |q_\phi - q_{\phi'}| \left(\prod_{j=1}^M (1 - q_j) \right).$

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Proof Approach.

Very weak bound on the utility of position $\hat{p} < p$.

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At most $\log_{1/\gamma^{**}}((q_1 - q_{M+1})/\epsilon)$ consecutive instances of Case 3 may occur.

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Proof Approach.

"Dynamical system in γ^{**} ." (Not dependent upon the form of $\gamma^{*}(q)$?)

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with Consumer Sea

Results

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Image: A mathematical states of the state

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• Convergence in an asynchronous bidding model

• Efficient convergence in an asynchronous bidding model



• Efficient convergence in an asynchronous bidding model

 Efficient convergence in an asynchronous bidding model with probability 1

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 - of the same form as that of Cary et al. (2008)

Road Map

• Position Auctions (without consumer search)

- Motivation
- Static Model (Edelman, Ostrovsky, and Schwarz (2007))
- Dynamic Model (Cary et al. (2008))
- Position Auctions with Consumer Search
 - Motivation
 - Static Model (Athey and Ellison (2008))
- Our Model
 - Convergence
- Generalizations?
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$\bullet\,$ Our method $\approx\,$ Cary et al. (2008)'s method

 Our method ≈ Cary et al. (2008)'s method; its applicability is naïvely surprising.

 $\bullet\,$ Our method $\approx\,$ Cary et al. (2008)'s method Three key steps:

• Our method pprox Cary et al. (2008)'s method

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- **o** proof that the *M* positions stabilize

restriction of the strategy space

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"Restricted Balanced Bidding"

restriction of the strategy space

- "Restricted Balanced Bidding"
- ensures that the equilibrium in each round is unique

analysis of low-quality advertisers's behavior

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our first lemma

analysis of low-quality advertisers's behavior

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analysis of low-quality advertisers's behavior

- our first lemma
- low-quality advertisers should drop out efficiently
- equilibrium bid monotonicity presumably unnecessary

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- Cases 1–3
- we showed: top *M* advertisers do not bid over their equilibrium bids "too often"
- may not hold in general—but is likely to hold when equilibrium is monotone

Conclusion

Image: A mathematical states of the state

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Conclusion

Convergence should be demonstrable in dynamic position auction models with sufficiently well-behaved static equilibrium strategies.

Questions?

Scott Duke Kominers (Harvard)

Dynamic Position Auctions

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